# **Maximum Flow and Minimum Cut**

- •Max flow and min cut.
  - Two very rich algorithmic problems.
  - Cornerstone problems in combinatorial optimization.
  - Beautiful mathematical duality.
- Nontrivial applications / reductions.
  - Data mining.
  - Project selection.
  - Airline scheduling.
  - Bipartite matching.
  - Image segmentation.
  - Network connectivity.
  - Chemical Production

- Network reliability.
- Distributed computing.
- Security of statistical data.
- Many many more . . .

## **Flow network**

- Abstraction for material flowing through the edges.
- G = (V, E) = directed graph, no parallel edges.
- Two distinguished nodes: s = source, t = sink.
- c(e) = capacity of edge e.



## Cuts

•Def. An s-t cut is a partition (A, B) of V with  $s \in A$  and  $t \in B$ .



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## **Minimum Cut Problem**

•Min s-t cut problem. Find an s-t cut of minimum capacity.



#### **Flows**

•Def. An s-t flow is a function that satisfies:

• For each  $e \in E$ : • For each  $v \in V - \{s, t\}$ :  $\sum_{e \text{ in to } v} f(e) = \sum_{e \text{ out of } v} f(e) \quad \text{(conservation)}$ • Def. The value of a flow f is:  $v(f) = \sum_{e \text{ out of } s} f(e) .$ 



#### **Flows**

• Def. An s-t flow is a function that satisfies:

• For each  $e \in E$ : • For each  $v \in V - \{s, t\}$ :  $\begin{array}{ll}
0 \leq f(e) \leq c(e) & \text{(capacity)} \\
f(e) = \sum f(e) & f(e) & \text{(conservation)}
\end{array}$ 

e in to v

•Def. The value of a flow f is:

$$f'(f) = \sum_{e \text{ out of } s} f(e)$$

e out of v



## **Maximum Flow Problem**

•Max flow problem. Find s-t flow of maximum value.



P1: Let f be any flow, and let (A, B) be any s-t cut. Then, the net flow sent across the cut is equal to the amount of flow leaving s.



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$$\sum_{e \text{ out of } A} f(e) - \sum_{e \text{ in to } A} f(e) = v(f)$$

#### Proof.

$$v(f) = \sum_{e \text{ out of } s} f(e)$$

By flow conservation, all terms except v = s are 0

$$= \sum_{v \in \mathcal{A}} \left( \sum_{e \text{ out of } v} f(e) - \sum_{e \text{ in to } v} f(e) \right)$$

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**P2:** Let f be any flow, and let (A, B) be any s-t cut. Then the value of the flow is at most the capacity of the cut.



**P2:** Let f be any flow. Then, for any s-t cut (A, B) we have  $v(f) \le cap(A, B)$ .



# **Certificate of Optimality**

- Max flow is at most equal to the capacity of the min cut (i.e., max flow is a lower bound to min cut)
- Let f be any flow, and let (A, B) be any cut.
  - If v(f) = cap(A, B), then f is a max flow and (A, B) is a min cut.

Value of flow = Cut capacity = 28  $\Rightarrow$  Max flow value = 28 = Min cut capacity

