

Ειδικά Θέματα Αλγορίθμων Ασκήσεις Φροντιστηρίου #4 RSA

1. The ciphertext 5859 was obtained from the RSA algorithm using $n = 11413$ and $e = 7467$. Using the factorization $11413 = 101 \cdot 113$, find the plaintext.
2. Let n be the product of two large primes. Alice wants to send a message m to Bob, where $\gcd(m, n) = 1$. Alice and Bob choose integers a and b , which are relatively primes to $\phi(n)$. Alice computes $c = m^a \pmod{n}$ and sends c to Bob. Bob computes $d = c^b \pmod{n}$ and sends d back to Alice. Since Alice knows a , she finds a_1 such that $aa_1 = 1 \pmod{\phi(n)}$. Then she computes $e = d^{a_1} \pmod{n}$ and sends e to Bob. Explain what Bob must now do to obtain m , and show that this works.
3. Naive Nelson uses RSA to receive a single ciphertext c , corresponding to the message m . His public modulus is n and his public encryption exponent is e . Since he feels guilty that his system was used only once, he agrees to decrypt any ciphertext that someone sends him, as long as it is not c , and return the answer to that person. Eve sends him the ciphertext $2^e c \pmod{n}$. Show how this allows Eve to find m .
4. Suppose Alice uses the RSA method as follows. She starts with a message consisting of several letters, and assigns $a = 1, b = 2, \dots, z = 26$. She then encrypts each letter separately. For example, if her message is cat, she calculates $3^e \pmod{n}$, $1^e \pmod{n}$ and $20^e \pmod{n}$. Then she sends the encrypted message to Bob. Explain how Eve can find the message without factoring n . In particular, suppose $n = 8881$ and $e = 13$. Eve intercepts the message:

4461 794 2015 2015 3603.

Find the message without factoring 8881.

5. Let $n = p \cdot q$ be the product of two distinct primes.
 - a) Let m be a multiple of $\phi(n)$. Show that if $\gcd(a, n) = 1$, then $a^m = 1 \pmod{p}$ and $1 \pmod{q}$.
 - b) Let m be a multiple of $\phi(n)$ and let a be arbitrary (possibly $\gcd(a, n) \neq 1$). Show that $a^{m+1} \equiv a \pmod{p}$ and $a \pmod{q}$.
 - c) Let e and d be encryption and decryption exponents for RSA with modulus n . Show that $a^{e \cdot d} \equiv a \pmod{n}$ for all a . This shows that we do not need to assume $\gcd(a, n) = 1$ in order to use RSA.
 - d) If p and q are large, why is it likely that $\gcd(a, n) = 1$ for a randomly chosen a ?