

Lecture 2

- Further Examples
- Dominance Relations



Further Examples

* Remember that:

- A metric space is a pair (X, d) where
 $X \neq \emptyset$, $d: X \times X \rightarrow \mathbb{R}$ such that:
 Carrier $\xrightarrow{\text{Metric}}$
 - i. $\forall x, y \in X, d(x, y) \geq 0$ (p.d.)
 - ii. $d(x, y) = 0 \Leftrightarrow x = y$ (sep.)
 - iii. $\forall x, y \in X, d(x, y) = d(y, x)$ (sym.)
 - iv. $\forall x, y, z \in X, d(x, y) \leq d(x, z) + d(z, y)$ (tf. ineq.)

- When only (\Leftarrow) in ii then \Rightarrow pseudo-metric
pseudo metric space

- Every $X \neq \emptyset$ can be equipped with the
discrete metric \Rightarrow the collection of metric spaces

is "real",

Exe: $c \in \mathbb{R}$, define $d_c(x, y) = \begin{cases} 0, & x=y \\ c, & x \neq y \end{cases}$ ✓

- For what c is d_c a (pseudo)-metric?
- Is d_c related to d_1 ?

- It is possible that the same carrier can be equipped with more than one metric. \square

Example : Hausdorff Distance - Information Theory

$n \in \mathbb{N}^*$, $X = \{0,1\}^n = \{(x_1, x_2, \dots, x_n), x_i \in \{0,1\}$ $\in \mathbb{R}^n$

$\forall i=1, \dots, n\}$ $\xrightarrow{\text{bits of information}}$ $\xrightarrow{\text{n-dimensional vectors}}$ $\xrightarrow{\text{0-1}}$ n -dimensional vectors

E.g. $n=2$, $X = \{(0,0), (0,1), (1,0), (1,1)\}$ etc

$d_H : X \times X \rightarrow \mathbb{R}$ is defined by :

$$\underline{x, y \in X}, d_H(x, y) := \# \{ \underline{i=1, \dots, n} : \underline{x_i \neq y_i} \}$$

$$\left. \begin{array}{l} n=2 \\ d_H((0), (1)) = 2 \\ d_H((0), (0)) = 0 \end{array} \right\}$$

"Number of disagreements between x, y "

$$- 0 \leq d_H(x, y) \leq n \quad \forall x, y \in X$$

hence d_H is a well defined real function and p.d. holds

- $\underbrace{D = d_H(x, y)} \Leftrightarrow x \text{ and } y \text{ Nowhere agree} \Leftrightarrow x = y$ [hence separation holds]
- $d_H(x, y) = \# \{i=1, \dots, n : x_i \neq y_i\} = \# \{i=1, \dots, n, y_i \neq x_i\}$
 $= d_H(y, x)$ [hence symmetry holds]
- if $x, y, z \in X$ then

$$\checkmark \# \left[\{i=1, \dots, n : x_i \neq z_i\} \cup \{i=1, \dots, n : z_i \neq y_i\} \right] \checkmark$$

→ consider the union

$$\leq \# \{i=1, \dots, n : x_i \neq z_i\} + \# \{i=1, \dots, n : z_i \neq y_i\} =$$

$$= \underline{d_H(x, z) + d_H(z, y)} \quad (\alpha)$$

And

→ since it can be that
 $x_i \neq z_i$ and $z_i \neq y_i$ but $x_i = y_i$

$$\begin{aligned} \{i=1, \dots, n : x_i \neq y_i\} &\subseteq \{i=1, \dots, n : x_i \neq z_i\} \cup \\ \{i=1, \dots, n : z_i \neq y_i\} &\Rightarrow \# \{i=1, \dots, n : x_i \neq y_i\} = \underline{d_H(x, y)} \end{aligned}$$

$$\leq \# \left[\{i=1, \dots, n : x_i \neq z_i\} \cup \{i=1, \dots, n : z_i \neq y_i\} \right], \quad (\beta)$$

$$(\alpha), (\beta) \Rightarrow \boxed{d_H(x, y) \leq d_H(x, z) + d_H(z, y)}$$

[hence tr. ineq. holds since x, y, z are arbitrary]

$\rightarrow d_A$ useful in error correction in message transmissions, etc. \square

Example: Metrics via P.d. Matrices

$n \in \mathbb{N}^*$, $X = \mathbb{R}^n$, A is an $n \times n$ strictly positive definite (symmetric) Matrix \Leftrightarrow

Positive definite (symmetric) Matrix \Leftrightarrow

$\forall x \in \mathbb{R}^n$ $x^T A x \geq 0$ with equality iff $x = 0_{n \times 1}$

\hookrightarrow column or row vector

\Leftrightarrow all the eigenvalues of A are strictly positive

real due to symmetry

e.g. $A = c \underbrace{\text{id}_{n \times n}}_{\text{c} > 0}, \text{id}_{n \times n} = \begin{pmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 1 \end{pmatrix}$

$d_A: \mathbb{R}^n \times \mathbb{R}^n \rightarrow \mathbb{R}$ is defined by

$$d_A(x, y) := \sqrt{(x-y)^T A (x-y)} \in \mathbb{R}^n$$

- Since n is finite and A is positive

$$\text{definite } 0 \leq d_A(x, y) < \infty \quad \forall x, y \in \mathbb{R}^n$$

Hence d_A is a well-defined real valued function and P.D. holds.

$$\begin{aligned}
 - \quad & \underline{0 = d_A(x,y)} \Leftrightarrow \underline{0 = ((x-y)' A (x-y))^{1/2}} \\
 \Leftrightarrow & \underline{\underline{0 = (x-y)' A (x-y)}} \Leftrightarrow \underline{\underline{x-y = 0_{n \times 1}}} \\
 & \text{A is strictly P.D.} \quad \Leftrightarrow \forall z \in \mathbb{R}^n \quad z' A z \geq 0 \\
 & \quad \quad \quad \text{with equality iff } z=0_n.
 \end{aligned}$$

$\Leftrightarrow x=y$ hence separation holds.

$$\begin{aligned}
 - \quad & \underline{d_A(y,x) = ((y-x)' A (y-x))^{1/2}} = \underline{\underline{((x-y)' A}} \\
 & \underline{\underline{[x,y])^{1/2}}} = \underline{\underline{(-1)^2 (x-y)' A (x-y))^{1/2}}} = \\
 & = \underline{\underline{(x-y)' A (x-y))^{1/2}}} = \underline{d_A(x,y)}
 \end{aligned}$$

Hence symmetry holds

- Exogenous: Cauchy-Schwarz inequality

$$\text{if } \underline{\underline{A}} \text{ is (P.D.) then } \underline{\underline{x^* A y^* \leq (x^* A x)^{1/2} (y^* A y)^{1/2}, \quad \forall x^*, y^* \in \mathbb{R}^n}}$$

Not necessarily
strictly positive definite.

Now if $x, y, z \in \mathbb{R}^n$ then

$$\underline{d_A^2(x,y)} = \underline{(x-y)'A(x-y)} = \underline{(x-z-y)'A(x-z-y)}$$

$$= \underline{(x-z) + (z-y)}' A \underline{(x-z) + (z-y)} =$$

$$= \underline{(x-z)' + (z-y)'} A \underline{(x-z) + (z-y)} =$$

$$= \underline{(x-z)'A(x-z)} + \underline{(x-z)'A(z-y)} + \underline{(z-y)'A(x-z)} +$$

$$\underline{(z-y)'A(z-y)} =$$

$$(\alpha)' = (\beta) \text{ but } (\alpha), (\beta) \text{ are } 1 \times 1 \Rightarrow (\alpha) = (\beta)$$

$$= \underline{(x-z)'A(x-z)} + \underline{(z-y)'A(z-y)} + 2\underline{(x-z)'A(z-y)}$$

$$= \underline{d_A^2(x,z)} + \underline{d_A^2(z,y)} + 2\underline{(x-z)'A(z-y)}$$

$$\leq \underline{d_A^2(x,z)} + \underline{d_A^2(z,y)} + 2 \left| \underline{(x-z)'A(z-y)} \right| \stackrel{\text{C.S. ineq.}}{\leq}$$

$$\leq \underline{d_A^2(x,y)} + \underline{d_A^2(z,y)} + 2 \left(\underline{(x-z)'A(x-z)} \right)^{1/2} \left(\underline{(z-y)'A(z-y)} \right)^{1/2}$$

$$x^* = x-z$$

$$y^* = z-y \quad d_A(x,z) \quad d_A(z,y)$$

$$= d_A^2(x, z) + d_A^2(z, y) + 2 \checkmark d_A(x, z) d_A(z, y) \checkmark$$

$$= (d_A(x, z) + d_A(z, y))^2 \checkmark$$

Hence we have shown:

$$d_A^2(x, y) \leq \underbrace{(d_A(x, z) + d_A(z, y))}_\text{Monotonicity}^2 \Rightarrow \checkmark$$

$d_A(x, y) \leq d_A(x, z) + d_A(z, y)$ [hence the triangle inequality holds]

* Show that when A is "only" positive definite then d_A is a pseudometric.

$$\begin{aligned} * \text{ When } A = \text{id}_{n \times n} \quad d_A(x, y) &= ((x-y)' I (x-y))^{1/2} \\ &= ((x-y)' (x-y))^{1/2} = \left((x_1 - y_1, \dots, x_n - y_n) \begin{pmatrix} x_1 - y_1 \\ \vdots \\ x_n - y_n \end{pmatrix} \right)^{1/2} \\ &= \left(\sum_{i=1}^n (x_i - y_i)^2 \right)^{1/2} = \text{Euclidean distance} \\ &\quad \text{sum of squared deviations} \end{aligned}$$

and (\mathbb{R}^n, d_I) is the "usual, Euclidean metric space.

* When $A = C \text{id}_{n \times n}$, $C > 0$ then

$$d_A(x, y) = ((x-y)' C (x-y))^{1/2} = \\ = \sqrt{C} ((x-y)' (x-y))^{1/2} = \sqrt{C} d_I(x, y)$$

[hence it is possible that different metrics on the same X - can obey relations between them]

* When $n=1$, $\mathbb{R}^n = \mathbb{R}$, $A = C$ with $C > 0$

and $d_A(x, y) = ((x-y)' C (x-y))^{1/2} =$

$$= \sqrt{C} \underline{(x-y)^2} = \sqrt{C} \underline{|x-y|} \quad \checkmark$$

When $C = \underline{\underline{1}}$ (and thereby $A = \text{id}_{1 \times 1}$)

$$\underline{d_I(x, y)} = |x-y| = \underline{d_U(x, y)}$$

And therefore the Euclidean metric is an "extension" of d_u on "higher dimensions". □

Example: Metric defined by the sum of absolute deviations

$X = \mathbb{R}^n$ (as before), consider

$\underline{d_{11}} : \mathbb{R}^n \times \mathbb{R}^n \rightarrow \mathbb{R}$ defined by
 $x, y \in \mathbb{R}^n$, $d_{11}(x, y) = \sum_{i=1}^n |x_i - y_i| =$
 $\underbrace{\sum_{i=1}^n d_u(x_i, y_i)}$

- Obviously $0 \leq d_{11}(x, y) < \infty \forall x, y \in \mathbb{R}^n$

hence d_{11} is a well defined real function and

p.d. holds.

- $0 = d_{11}(x, y) \Leftrightarrow 0 = \sum_{i=1}^n |x_i - y_i| \Leftrightarrow$
 $|x_i - y_i| = 0 \quad \forall i = 1, \dots, n \Leftrightarrow x_i = y_i \quad \forall i = 1, \dots, n \Leftrightarrow x = y$

Hence separation holds.

$$\forall d_u(x_i, y_i) = 0 \quad \forall i = 1, \dots, n$$

$$\begin{aligned}
 - d_{11}(y, x) &= \sum_{i=1}^n |y_i - x_i| \stackrel{\substack{= |x_i - y_i| \text{ if } i=1, \dots, n}}{=} \sum_{i=1}^n |x_i - y_i| = \\
 &= d_{11}(x, y) \text{ hence symmetry holds}
 \end{aligned}$$

$$\begin{aligned}
 - \text{ if } x, y, z \in \mathbb{R}^n, d_{11}(x, y) &= \sum_{i=1}^n |x_i - y_i| \\
 &= \sum_{i=1}^n |x_i - z_i - y_i| = \sum_{i=1}^n |(x_i - z_i) + (z_i - y_i)| \\
 &\leq \sum_{i=1}^n (|x_i - z_i| + |z_i - y_i|) = \sum_{i=1}^n |x_i - z_i| + \sum_{i=1}^n |z_i - y_i| = \\
 &= \underline{d_{11}(x, z) + d_{11}(z, y)}
 \end{aligned}$$

Once again
 we "interpolate",
 use triangle ineq.
 for absolute value

\leq
 tr. ineq.
 for absolute
 value

hence the triangle inequality holds

* When $n=1$ $d_{11}(x, y) = |x - y| = d_{\text{eu}}(x, y)$
 hence d_{11} is also our extension of due to
 "higher dimensions".

$$\begin{aligned}
 * \text{ Eg. } n=2 \quad d_{\parallel}((\overset{2}{1}, \overset{1}{2})) &= \\
 &= |2-1| + |1-2| = \boxed{2} \\
 d_{\perp}((\overset{2}{1}, \overset{1}{2})) &= \sqrt{(\overset{2}{1}-\overset{1}{1})^2 + (\overset{1}{2}-\overset{2}{1})^2} = \\
 &= \boxed{\sqrt{2}}
 \end{aligned}$$

Hence $d_{\perp} \neq d_{\parallel}$.

□

Example: Metric defined via the Maximum absolute deviation

$$\underline{X = \mathbb{R}^n} \text{ (as before)}$$

$d_{\max} : \mathbb{R}^n \times \mathbb{R}^n \rightarrow \mathbb{R}$ defined by

$$x, y \in \mathbb{R}^n, d_{\max}(x, y) := \max_{i=1, \dots, n} |x_i - y_i| = \max_{i=1, \dots, n} d_{\parallel}(x_i, y_i)$$

- Since n finite, $0 \leq \max_{i=1, \dots, n} |x_i - y_i| < +\infty$

Hence d_{\max} is a well defined real function

that also satisfies p.d.

$\max_{i=1,\dots,n}$

- $D = d_{\max}(x, y) \Leftrightarrow D = \max_i |x_i - y_i| \Leftrightarrow$
 $|x_i - y_i| = 0 \quad \forall i=1,\dots,n \Leftrightarrow x_i = y_i \quad \forall i=1,\dots,n$ hence separation holds. \rightarrow abs value symmetric
- $d_{\max}(y, x) = \max_i |y_i - x_i| = \max_i |x_i - y_i|$
 $= d_{\max}(x, y)$ hence symmetry holds

- if $x, y, z \in \mathbb{R}^n$, $d_{\max}(x, y) = \max_{i=1,\dots,n} |x_i - y_i|$

$$= \max_{i=1,\dots,n} |x_i + z_i - y_i| \leq \max_{i=1,\dots,n} (|x_i - z_i| + |z_i - y_i|)$$

remember the properties of sup + tr. ineq. for the abs. value

$$\leq \max_{i=1,\dots,n} |x_i - z_i| + \max_{i=1,\dots,n} |z_i - y_i| = d_{\max}(x, z) + d_{\max}(z, y)$$

Again property of sup hence the triangle inequality holds.

* When $n=1$, $d_{\max}(x, y) = \max|x-y| = |x-y|$
 $= d_u(x, y)$ hence d_{\max} is yet another extension of d_u

$$* n=2, d_{\max}((\underline{\underline{z}}_1), \underline{\underline{z}}_2) =$$

$$= \max\{12-11, 11-2\} = \max\{111, 111\}$$

$$= \underline{\underline{1}} \neq d_{11}((\underline{\underline{z}}_1), (\underline{\underline{z}}_2))^{=2}$$

$$\times d_I((\underline{\underline{z}}_1), (\underline{\underline{z}}_2))^{=2} \quad \text{||}$$

More generally for $n > 1$

$$(\mathbb{R}^n, d_I) \neq (\mathbb{R}, d_{11})$$

~~X~~ ~~#~~ ✓

$$(\mathbb{R}^n, d_{\max})$$

However notice: $\forall x, y \in \mathbb{R}^n$

$$(a) - d_{\max}(x, y) = \max |x_i - y_i| \leq \sum_{i=1}^n |x_i - y_i| = d_{11}(x, y)$$

$$(b) - d_{11}(x, y) = \sum_{i=1}^n |x_i - y_i| \leq \sum_{i=1}^n \max_{i=1, \dots, n} |x_i - y_i|$$

$$= n \cdot \max_{i=1, \dots, n} |x_i - y_i| = n d_{\max}(x, y)$$

$$(c) - d_{\max}^2(x, y) = (\max |x_i - y_i|)^2 = \max (|x_i - y_i|^2)$$

$$\leq \sum_{i=1}^n (x_i - y_i)^2 = d_I^2(x, y) \Rightarrow d_{\max}(x, y) \leq d_I(x, y)$$

monotonicity of r .

$$(S.) - d_I^2(x, y) = \sum_{i=1}^n (x_i - y_i)^2 \leq \sum_{i=1}^n \max(x_i - y_i)^2$$

$$\leq n \cdot \max(x_i - y_i)^2 = n (\max|x_i - y_i|)^2$$

$$= n d_{\max}(x, y) \Rightarrow d_I(x, y) \leq \sqrt{n} d_{\max}(x, y)$$

Non.
of $\sqrt{\cdot}$.

Hence by (a), (b)

$$d_{\max} \leq d_{II}$$

$$< d_{II} \leq n d_{\max}$$

by (g), (g)

$$d_{\max} \leq d_I$$

$$d_I \leq \sqrt{n} d_{\max}$$

$d_{\max} \leq d_I \leq n d_{\max}$ by (b), (g)

$$d_{II} \leq n d_I$$

by (a), (g)

$$d_I \leq \sqrt{n} d_{II}$$

$$\frac{1}{n} d_{II} \leq d_I \leq \sqrt{n} d_{II}$$

$d_{\max} \leq d_{II} \leq n d_{\max}$

(functional inequality
hold for x, y)

Hence the three metrics obey functional relations!

* More generally for d_1, d_2 metrics defined on X :

i. d_2 dominates d_1 iff ✓

$\exists c > 0$: $d_1 \leq c d_2$ [functional inequality
 $\Leftrightarrow d_1(x, y) \leq c d_2(x, y) \forall x, y \in X$]

ii. d_1 is equivalent to d_2 iff ✓

d_2 dominates d_1 and d_1 dominates d_2

Exercise: d_1 is equivalent to d_2 iff $\exists c^*, C^* > 0$

$$: \underbrace{c^* d_2 \leq d_1 \leq C^* d_2}_{}$$

Hence d_I , d_{II} , and d_{max} are pairwise equivalent

* i.e. different metrics on the same X can obey dominance relations [based on functional inequalities]

→ We will later on the course see what this implies in terms of properties ↗

End of lecture 2.

Example : function Space

For $V \neq \emptyset$ consider $BC(V, \mathbb{R})$ and let

$d_{\text{sup}} : BC(X, \mathbb{R}) \times BC(V, \mathbb{R}) \rightarrow \mathbb{R}$ be defined

by: $f, g \in BC(X, \mathbb{R})$, $d_{\text{sup}}(f, g) := \sup_{x \in V} |f(x) - g(x)|$

Notice that

$$\begin{aligned} * \quad & \sup_{x \in V} |f(x) - g(x)| \leq \sup_{x \in V} (|f(x)| + |g(x)|) \\ & \leq \sup_{x \in V} |f(x)| + \sup_{x \in V} |g(x)| < +\infty \\ & f, g \in BC(V, \mathbb{R}) \end{aligned}$$

Hence d_{sup} is a well defined real function.

* Obviously $\sup_{x \in V} |f(x) - g(x)| \geq 0 \quad \forall f, g$
hence p.d. holds

* $D = d_{\text{sup}}(f, g) \Leftrightarrow D = \sup_{x \in V} |f(x) - g(x)|$
 $\Leftrightarrow |f(x) - g(x)| = D \quad \forall x \in V \Leftrightarrow f(x) = g(x) \quad \forall x \in V$

$\Leftarrow f=g$ hence separation holds

$$* \text{dsep}(g,f) = \sup_{x \in V} |g(x) - f(x)| = \sup_{x \in V} |f(x) - g(x)|$$

symmetry
of $| \cdot |$

$$= \text{dsep}(f,g), \text{ hence symmetry holds.}$$

* if $f, g, h \in BC_V(\mathbb{R})$

$$\text{dsep}(f,g) = \sup_{x \in V} |f(x) - g(x)| = \sup_{x \in V} |f(x) + h(x) - g(x)|$$

$$= \sup_{x \in V} |(f(x) - h(x)) - (h(x) - g(x))| \leq \sup_{x \in V} |f(x) - h(x)| +$$

$$+ |h(x) - g(x)| \leq \sup_{x \in V} |f(x) - h(x)| + \sup_{x \in V} |h(x) - g(x)|$$

$= \text{dsep}(f,h) + \text{dsep}(h,g)$ hence the triangle inequality holds.

To BE CONTINUED ...