OIKONOMIKO ΠΑΝΕΠΙΣΤΗΜΙΟ ΑΘΗΝΩΝ



ATHENS UNIVERSITY
OF ECONOMICS
AND BUSINESS

Information-Centric Networks

Section # 4.1: Routing Issues

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Week 4 / Paper 1

- Open issues in Interdomain Routing: a survey
 - Marcelo Yannuzzi, Xavier Masip-Bruin, Olivier Bonaventure
 - IEEE Network, Nov.-Dec. 2005, vol. 19, no. 6, pp. 49 56
- Main point
 - There are many challenges in interdomain routing
 - Relationships between challenges
 - Review of problems and proposed solutions
 - Reasons for non-adoption of solutions

Introduction

- Limitations of BGP
 - Hard to replace due to widespread deployment
 - Growth of Internet strains the protocol
- Routing among distinct domains
 - Independent management of domains
 - Active competition between providers
- Basics of interdomain routing
 - Autonomous System (AS): a distinct network on the Internet
 - Managed by a single authority
 - Common internal routing policy
 - Intradomain routing via IS-IS or OSPF
 - More than 20000 ASes in the Internet

Basics of interdomain routing

- Types of ASes
 - Single homed stub ASes
 - The "leaves" of the network
 - Multi homed stub ASes
 - Load balancing or failure resiliency
 - Transit ASes
- Types of AS relationships
 - Costumer-provider
 - Peer-peer
- The AS hierarchy
 - Tier-1 (no upstream provider) composing the internet core
 - Tier-2 are customers of Tier-1 and providers of Tier-2
 - ASes also use peering links to directly exchange traffic

Basics of interdomain routing

- Interior and exterior BGP
 - BGP routers exchange reachability information via eBGP
 - This information is distributed internally via iBGP
- Route discovery and selection
 - BGP routers advertise AS level paths to IP prefixes
 - There is no global topology view
 - The path is abstracted into AS numbers
 - Route selection is heavily policy influenced
 - 1. Choose the route with the highest local preference
 - 2. Choose the route with the shortest AS path
 - Choose the route with the lowest Multi Exit Discriminator
 - 4. Prefer the route with the lowest IGP metric
 - 5. Run tie-breaking rules

- Limited traffic engineering capabilities
 - Only reachability is advertised
 - No way to propagate multiple routes
 - Inbound traffic is hard to control with BGP
 - It is the sender and not the receiver that chooses paths
 - This is due to the uncoordinated routing on the Internet
- Lack of QoS support
 - Many ASes have deployed differentiated services
 - These allow coarse grained intradomain QoS
 - But BGP does not allow interdomain QoS
 - Most ASes prefer over-provisioning instead of QoS

- How has the Internet changed?
 - AS numbers have swelled
 - Connections per AS have increased
 - Additional applications
 - Application requirements are not reflected into BGP
- Slow convergence and chattiness
 - BGP messages: Open, Update, Notification, Keepalive
 - A failure causes large amounts of BGP updates
 - Path exploration takes place until convergence
 - BGP routers wait for MRAI before updating a route
 - BGP routers often employ route flap damping
 - Ignore routes that change too often
 - Improve stability at the expense of convergence

- Slow convergence and chattiness
 - Proposal: faster propagation of updates due to failures
 - Two methods, ghost flushing and reporting the root cause
 - Limits path exploration
 - Requires BGP modification to indicate a failure cause
 - Very hard to pinpoint a failure due to route aggregation
 - Route disaggregation impacts scalability, so there is a tradeoff
 - Proposal: infer source of failures by correlating data
 - Requires multiple vantage points and offline processing
 - Each solution adversely affects some objectives
 - Solutions: root-cause, MRAI timer, flap damping, aggregation
 - Objectives: scalability, convergence, message load

- Scalability problems due to multihoming
 - Many stub ASes are multihomed
 - Resilience and load balancing
 - Each such AS has multiple IP prefixes (one per provider)
 - Each prefix is advertised to all providers
 - Different paths are advertised to indicate route preference
 - Depending on provider aggregation rules, there may be problems
 - Eventually, disaggregation may be required to achieve policy goals
 - Also, each provider can only aggregate its own prefixes
 - All these lead to even more routes being advertised
 - Proposal: route filtering
 - Avoid propagating very long prefixes (very specific routes)
 - This inhibits load balancing (it hides some routes)
 - What we really need is better support for multihoming

- Expresiveness and safety of policies
 - Each domain independently chooses its policies
 - The result is suboptimal due to lack of coordination
 - Global routing anomalies arise
 - Global divergence of routing policies
 - BGP policies are not that expressive
 - Rich enough to express intricate routing policies
 - Not rich enough to allow discovery of problems
 - Each AS does not want to disclose its internal details
 - Many problems are hidden inside the AS's network

- Robustness of BGP sessions
 - BGP routers communicate via TCP
 - Network congestion can lead to failures
 - Hard to resolve congestion due to routing problems
 - Need to distinguish routing messages from ordinary traffic
- Security issues
 - Spoofed TCP RSTs can bring down BGP sessions
 - Also spoofed TCP messages can be inserted
 - Filter spoofed packets
 - Use authentication between BGP routers
 - BGP advertisements are not authenticated
 - S-BGP certifies the validity of routes by signatures
 - Processing cost and need for a PKI

- Lack of multipath routing
 - BGP routers only advertise a single path per prefix
 - Even if they have received many alternatives
 - BGP routers may only use a single path per prefix
 - · Some implementations use many for load balancing
 - Proposed extensions raise scalability concerns
- Transit through an AS: iBGP issues
 - A large AS has trouble propagating routes inside it
 - Ideally all internal routers should communicate with each other
 - This not scalable for large ASes
 - Different routers may treat traffic differently
 - Encapsulation: guarantees that packets will use preferred routes
 - Operates between ingress and egress BGP routers

End of Section #4.1

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