ΟΙΚΟΝΟΜΙΚΟ ΠΑΝΕΠΙΣΤΗΜΙΟ ΑΘΗΝΩΝ



ATHENS UNIVERSITY OF ECONOMICS AND BUSINESS

### M.Sc. Program in Data Science Department of Informatics

### **Optimization Techniques Convex Optimization**

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# Outline

- Convex sets
  - Definitions and basic concepts
- Convex functions
  - Equivalent definitions
  - Advantages when optimizing convex functions
- Convex optimization problems
  - Unconstrained optimization
    - Descent methods
  - Constrained optimization
    - Lagrange duality and the KKT conditions
    - Algorithms

# Our goals

- Formulate problems where the objective function or the constraints are not linear
- Understand when can we have efficient algorithms for solving "non-linear" programs
  - What assumptions are needed for the type of constraints or for the objective function?
- Generalize LP duality theory/sensitivity analysis?

# Introduction to convex sets and convex functions

We focus on subsets of  $\mathbb{R}^n$  for some dimension  $n \ge 1$ 

- Points here correspond to n-dimensional vectors
- But intuition from low dimensions very useful
- Given 2 points x, y ∈ R<sup>n</sup>, a point z lies on the line that connects x and y if and only if

z = ax + (1-a) y for some  $a \in [0, 1]$ 

**Definition:** A set  $C \subseteq \mathbb{R}^n$  is convex if for any 2 points  $x, y \in C$ , and for any  $a \in [0, 1]$ , we have that  $ax + (1-a)y \in C$ 

 In geometric terms: the line connecting any 2 points of C, must entirely belong to C



#### **Examples**

#### Convex sets:



Nonconvex sets:



#### **Further examples of convex sets**

1.All of  $\mathbb{R}^n$ , for any dimension  $n \ge 1$ 

2. The nonnegative orthant: points with all coordinates nonnegative

• Since ax + (1-a)y will also have nonnegative coordinates

3. The set of points contained within a ball

- E.g.  $\{x: ||x||_2 \le 1\}$
- Since

 $||ax + (1-a)y||_2 \le ||ax||_2 + ||(1-a)y||_2 \le a||x||_2 + (1-a)||y||_2 \le 1$ 

4. Intersections of convex sets

#### **Further examples of convex sets**

- 5. Feasible region of a linear program
  - Convex polygon in 2 dimensions (when it is bounded)
  - Convexity follows since it is an intersection of halfspaces



#### **Examples that do not involve R**<sup>n</sup>

 The exact same definition of convexity can be applied for elements that are not points of R<sup>n</sup>

**Definition:** A real symmetric n x n matrix A is called positive semidefinite (PSD) if for every n-dimensional vector z  $z^{T} \cdot A \cdot z \ge 0$ 

Claim: The set of PSD matrices is a convex set i.e., if A and B are PSD matrices, then  $\lambda A + (1-\lambda)B$  is also PSD for any  $\lambda \in [0, 1]$ 

**Definition:** A function  $f : \mathbb{R}^n \to \mathbb{R}$  is convex if for any 2 points x, y, and for any  $a \in [0, 1]$ , we have that  $f(ax + (1-a)y) \le af(x) + (1-a)f(y)$ 



 Geometric interpretation: the line connecting any 2 points (x,f(x)) (y,f(y)) must lie on or above the graph of the function

#### **Examples**

- •With 1 variable:
  - Exponential functions with base > 1:  $2^x$ ,  $e^x$ ,  $c^x$  for  $c \ge 1$
  - Polynomial functions:  $x^3$ ,  $x^{10}$ ,  $x^c$ , for  $c \ge 1$
  - Linear functions: we have exact equality in the definition
- •With many variables
  - Exponential functions:  $e^{x+y}$ ,  $e^{x+y+z}$ ,
  - Negative of logarithms: log(x + y)
  - The sum of convex functions remains convex

#### **Equivalent definitions**

(1)Based on the epigraph

The epigraph of a function f is the set:

epi f = { (x, t): 
$$t \ge f(x)$$
 }

f is convex if and only if the epigraph of f is a convex set

• Geometric interpretation: the set of points that lie on or above the graph of the function should be a convex set



#### **Equivalent definitions**

(2) Based on the partial derivatives

Suppose that f is twice differentiable

For functions with 1 variable: f is convex if and only if  $f''(x) \ge 0$ , for every x

•The first derivative is increasing



e.g. for  $f(x) = x^2$ , f''(x) = 2 for every x

#### **Equivalent definitions**

#### (2) Based on the partial derivatives

For functions with n variables: Define the Hessian of f at point x as the n x n array:

$$H(f,x) = \begin{bmatrix} \frac{\partial^2 f(x)}{\partial^2 x_1} & \frac{\partial^2 f(x)}{\partial x_1 \partial x_2} & \cdots & \frac{\partial^2 f(x)}{\partial x_1 \partial x_n} \\ \frac{\partial^2 f(x)}{\partial x_2 \partial x_1} & \frac{\partial^2 f(x)}{\partial^2 x_2} & \cdots & \frac{\partial^2 f(x)}{\partial x_2 \partial x_n} \\ \vdots & \vdots & \ddots & \vdots \\ \frac{\partial^2 f(x)}{\partial x_n \partial x_1} & \frac{\partial^2 f(x)}{\partial x_n \partial x_2} & \cdots & \frac{\partial^2 f(x)}{\partial^2 x_n} \end{bmatrix}$$

A function f is convex if and only if the Hessian is positive semidefinite for every x

• Easy to see from Slide 13 that this holds for 1 dimension

#### **Equivalent definitions**

(3) Based on the tangents to the graph of f

f is convex if and only if the graph of f lies on or above all its tangents:



#### Algebraically in 1 dimension:

Slope of tangent at x = the derivative of f at x
Hence, if f lies above all its tangents, then for every x, y:

 $f(y) \ge f(x) + f'(x) (y-x)$  (\*)

#### **Equivalent definitions**

(3) Based on the tangents to the graph of f

#### For functions with n variables:

• Recall the gradient of a function

$$\nabla f(x) = \left(\frac{\partial f(x)}{\partial x_1}, \frac{\partial f(x)}{\partial x_2}, \dots, \frac{\partial f(x)}{\partial x_n}\right)$$

- The gradient shows the rate of increase/decrease along each dimension just as the derivative does for one variable
- Generalizing (\*) for many variables:

 $f(y) \ge f(x) + \nabla f(x)^{\top} \cdot (y-x) \quad (**)$ 

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One of the most important properties of convex functions

Sometimes we may also discuss concave functions

**Definition:** A function  $f : \mathbb{R}^n \to \mathbb{R}$  is concave if for any 2 points x, y, and for any  $a \in [0, 1]$ , we have that  $f(ax + (1-a)y) \ge af(x) + (1-a)f(y)$ 

- If f is concave, -f is convex
- Hence, maximizing a concave function f can be reduced to minimizing a convex function

### **Convex Optimization Problems**

# **Nonlinear Optimization Problems**

General form of optimization problems:

• Both equality and inequality constraints present

min f(x)s. t.:  $g_i(x) \le 0, \quad i = 1, 2, ..., m$  $h_i(x) = 0, \quad i = 1, 2, ..., p$ 

We say the above is a convex optimization problem when

- f(x) is a convex function
- Each g<sub>i</sub> is a convex function
- Each  $h_i$  is an affine function,  $h_i = a_i^T x b_i$

# **Convex Optimization**

Applications of convex optimization:

- Machine learning: linear regression (least squares), classification (logistic regression, support vector machines)
- Statistics: parameter estimation
- Control theory
- Signal processing
- And many many more...

# **Convex Optimization**

- For general non-convex problems, almost no hope
- There is no general approach that can work for any arbitrary optimization problem
- Some families of non-convex problems can be handled
- But when working under assumptions like convexity, or related properties (e.g. strong convexity), we can have guarantees for convergence and running time
- Still however not a standard technology, contrary to LP solvers
  - Commercial availability not as large as for LP solving but gradually changing

Unconstrained Convex Optimization

- We start with the easier version that has no constraints
- Suppose we just want to minimize a function f : R<sup>n</sup> → R without any further constraints
  - Still interesting problem with many applications
- Assumption: f is twice continuously differentiable
- Necessary condition for a point  $x^*$  to be a minimum is  $\nabla f(x^*) = 0$
- **BUT**: for an arbitrary function f:
  - This is not a sufficient condition, many other points may satisfy this (such as local optima)

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- We start with the easier version that has no constraints
- Suppose we just want to minimize a function  $f: \mathbb{R}^n \to \mathbb{R}$  without any further constraints
  - Still interesting problem with many applications
- Assumptions from now on (unless otherwise stated)
  - f is **convex**, and twice continuously differentiable
  - The minimum of f is attained (and  $\neq +\infty$  or  $-\infty$ )

Why is it nice to be convex:

Theorem: For a convex function f, a point  $x^*$  is a global minimum of f if and only if  $\nabla f(x^*) = 0$ 

Proof:

Recall the basic property of convex functions, i.e., inequality (\*\*):  $f(y) \ge f(x) + \nabla f(x)^T \cdot (y-x)$  for any 2 points x, y

- Suppose there exists  $x^*$  for which  $\nabla f(x^*) = 0$
- Then for every point y, inequality (\*\*) implies  $f(y) \ge f(x^*)$
- Hence x<sup>\*</sup> is a global minimum

#### Why is it nice to be convex:

- The theorem makes our lives much easier (not trivial however)
  - It suffices to find a point where the derivatives become 0
  - Local minima are global minima, as with linear programs (recall the terminating condition of simplex)
  - If we can solve analytically the system  $\nabla f(x) = 0$ , then no need for an algorithm
- In many cases convexity helps us exploit the geometric intuition we have from polyhedra or linear programming problems

#### Algorithms for convex unconstrained optimization:

- Iterative algorithms, updating a current feasible solution
- They produce a sequence of points x<sup>(0)</sup>, x<sup>(1)</sup>,..., x<sup>(k)</sup> with the property that

 $f(x^{(k)}) \rightarrow p^*$  as  $k \rightarrow \infty$ 

- $p^*$  is what we are after:  $p^* = inf_x f(x)$
- We may never find the actual optimal solution
- But we can get very close, in fact arbitrarily close if we allow enough iterations
- We can view these algorithms as iterative methods for solving the system ∇f(x) = 0

### **Descent Methods**

General form of descent methods

- Make a local update towards an appropriate direction
- Stop when  $\nabla f(x)$  is close to 0
- Initialization: k=0, pick a starting point  $x^{(0)}$ , and a step size  $\alpha_0$
- Update:
  - Check if stopping criterion satisfied
  - If not,  $x^{(k+1)} = x^{(k)} + \alpha_k \Delta x^{(k)}$
  - k++
- Usual stopping criterion:  $\|\nabla f(\mathbf{x}^{(k)})\|_2 \leq \epsilon$

Terminology:

- • $\Delta x$ : search direction
- • $\alpha_k$ : step size, with  $\alpha_k > 0$

### **Descent Methods**

How should we pick the search direction?

- Need to ensure that for every iteration k,  $f(x^{(k+1)}) \le f(x^{(k)})$
- Convexity, i.e. using (\*\*), implies it suffices to enforce that:  $\nabla f(x^{(k)})^T \cdot \Delta x^{(k)} < 0$
- Hence: choosing the (negative) gradient itself for the search direction is a safe choice!

# **The Gradient Descent Method**

One of the simplest algorithms in optimization: Descend according to the gradient direction

- Initialization: k=0, pick a starting point  $x^{(0)}$ , and a step size  $\alpha_0$
- Update: –Check if stopping criterion satisfied –If not,  $x^{(k+1)} = x^{(k)} - \alpha_k \nabla f(x^{(k)})$ –k++
- Stopping criterion  $\|\nabla f(x^{(k)})\|_2 \leq \epsilon$ 
  - If e.g.,  $\|\nabla f(x^{(k)})\|_2 \le (2m\epsilon)^{1/2}$ , where m is a lower bound on the minimum eigenvalue of H(f,x), then f(x) p<sup>\*</sup>  $\le \epsilon$

### **The Gradient Descent Method**

How should we pick the step size  $\alpha_k$ ?

- First idea: Exact line search
  - Find the minimum value of f along the gradient direction:

 $\alpha_k = \operatorname{argmin}_s f(x^{(k)} - s \nabla f(x^{(k)}))$ 

- 1-dimensional problem
- E.g., we could solve it via Newton's method
- But often too time consuming in practice

## **The Gradient Descent Method**

How should we pick the step size  $\alpha_k$ ?

- Second idea: Backtracking line search, an approximate solution to the exact line search
  - Try to approximately minimize f along the ray x  $s\nabla f(x^{(k)})$
  - Essentially make sure the function decreases "enough"
  - Many variants in the literature, e.g.

Keep setting s:= βs until

 $\begin{aligned} f(x - s \nabla f(x^{(k)})) &\leq f(x^{(k)}) - \alpha s \cdot ||\nabla f(x^{(k)}) ||_2^2 \\ \text{for } \beta < 1, \alpha < 1/2 \end{aligned}$ 

– Works well in practice

### **Descent Methods**

Example 1: Consider the function  $f(x_1, x_2) = x_1^2 + 2x_2^2 - 2x_1x_2$ 

Execute the first 2 steps of gradient descent with exact line search, starting from  $x^{(0)} = (1, 1)$ 

#### **Descent Methods**

#### Example 2: Consider the function $f(x_1, x_2) = \frac{1}{2}(x_1^2 + \gamma x_2^2), \quad \gamma > 0$

Start at  $x^{(0)} = (\gamma, 1)$ After k iterations of gradient descent, we get:

$$x^{(k)} = \left(\gamma\left(\frac{\gamma-1}{\gamma+1}\right)^k, \left(-\frac{\gamma-1}{\gamma+1}\right)^k\right)$$

Run with  $\gamma = 10$ 



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# **Convergence analysis**

Can we establish convergence properties for the gradient descent method?

- Empirically, it works well on average for convex functions
- Theoretically, upper bounds can be obtained when assuming strong convexity

Definition: A function is strongly convex when there exists m>0 such that for any x,

 $H(f, x) \ge m \cdot I$ 

- I is the identity matrix

# **Convergence analysis**

- Strong convexity together with (\*\*) implies that there also exists upper bounds on the Hessian
- Hence, there exist m>0 and M>0 such that for every x:

#### $m \cdot l \le H(f, x) \le M \cdot l$

- Convergence results on the number of iterations depend on
  - m and M
  - The initial solution  $x^{(0)}$
  - The accuracy parameter in the stopping criterion
- Note: we may not be aware of the values for m and M
  - It might be difficult to estimate for some functions
  - So, we may not know how many iterations we need
- Still, these bounds are conceptually useful
  - They provide a guarantee that the method converges

# **Convergence** analysis

- A relatively loose analysis with exact line search
- Theorem: For strongly convex functions, the number of iterations required by the gradient descent method is bounded by

 $\log((f(x^{(0)}) - p^*)/\epsilon) / \log(1/c)$ 

where

- c = 1 m/M < 1
- $p^* = min_x f(x)$
- $-\epsilon$  = accuracy parameter (= final suboptimality)
- $f(x^{(0)}) p^* = initial suboptimality$
- Thus, nominator = log of initial suboptimality to final suboptimality
- Conclusions: The error  $f(x^{(k)}) p^*$  converges to 0 at least as fast as a geometric series
  - i.e., linear convergence
- With backtracking line search, slightly worse bounds can also be established

A different descent method with favorable performance

- It is instructive to see first the method in 1 dimension
  - When n=1, we search for a point x, where f'(x) = 0
  - Suppose after k iterations, we have reached a point  $x_k$
  - How shall we move to the next iteration and pick  $x_{k+1}$ ?

#### Newton's method for n=1:

$$\mathbf{x}_{k+1} = \mathbf{x}_k - (\mathbf{f}'(\mathbf{x}_k) / \mathbf{f}''(\mathbf{x}_k))$$

Also referred to as the Newton-Raphson method

#### Newton's method for n=1:

$$x_{k+1} = x_k - (f'(x_k)/f''(x_k))$$



#### Geometric interpretation:

•Consider the plot of the derivative f'

•By convexity the first derivative is an increasing function

• Draw the tangent at x<sub>k</sub>

•Slope of the tangent =  $f''(x_k)$ 

•Find the point where the tangent hits the x-axis

•This is given by solving the equation

 $0 = f'(x_k) + f''(x_k)(x - x_k)$ 

Newton's method for n=1:

$$x_{k+1} = x_k - (f'(x_k)/f''(x_k))$$



#### Algebraic intuition:

• Consider the 2<sup>nd</sup> order Taylor approximation:

$$f(x_{k+1}) = f(x_k) + f'(x_k)(x_{k+1} - x_k) + f''(x_k)(x_{k+1} - x_k)^2/2$$

- •How would we choose to move from x<sub>k</sub> to  $x_{k+1}$ ?
- •Set derivative (with respect to  $x_{k+1}$ ) = 0  $\Rightarrow$  x<sub>k+1</sub> = x<sub>k</sub> - f'(x)/f''(x)
- • $X_{k+1}$  is the minimizer of g

• If f is close to a quadratic function, then the Newton step is close to the best possible

For many variables, we can generalize the same intuition:

- 2<sup>nd</sup> order Taylor approximation for a function of n variables
- Now x and  $\delta$  are n-dimensional vectors

 $f(\mathbf{x} + \boldsymbol{\delta}) = f(\mathbf{x}) + \nabla f(\mathbf{x})^{\mathsf{T}} \cdot \boldsymbol{\delta} + \frac{1}{2} \boldsymbol{\delta}^{\mathsf{T}} \cdot \mathsf{H}(f, \mathbf{x}) \cdot \boldsymbol{\delta}$ 

If we try to minimize with respect to  $\delta$  (= $\Delta x$ ), we get that:  $\delta$  = - H(f, x)<sup>-1</sup>  $\nabla f(x)$ 

- Is this a descent direction?
- To be aligned with the convexity of f we need to check that:  $-\nabla f(x)^{T} \cdot (H(f, x)^{-1} \cdot \nabla f(x)) < 0,$

But the Hessian is a PSD matrix!

Summarizing:

- Initialization: k=0, pick a starting point  $x^{(0)}$ , and a step size  $\alpha_0$
- Update:
  - Check if stopping criterion satisfied
  - If not,  $x^{(k+1)} = x^{(k)} \alpha_k H(f, x^{(k)})^{-1} \cdot \nabla f(x^{(k)})$
  - k++
- Usual stopping criterion:
  - Let  $\lambda := (\nabla f(x^{(k)})^{\mathsf{T}} \cdot H(f, x^{(k)}) \cdot \nabla f(x^{(k)}))^{1/2}$
  - Stop when  $1/2\lambda^2 \le \epsilon$
  - $\lambda$  is called the Newton decrement
  - Useful parameter for the analysis of the method

• Progress made using the 2<sup>nd</sup> order approximation



#### • Pros

- It is fast in general
- Scales well with problem size
- Performance not depend on problem parameters (?)
- Cons
  - Cost of computing the Hessian
- Convergence analysis
  - Can be established in a similar way as with gradient descent
  - Theoretical upper bound: proportional to  $f(x^{(0)}) p^*$