

# Entropy Exercise

Multimedia Technology

Tutorial 1, section 2

# Entropy Calculation Exercise

Assume an information source that produces eight symbols with probabilities  $1/4$ ,  $1/4$ ,  $1/8$ ,  $1/8$ ,  $1/8$ ,  $1/16$ ,  $1/32$ , and  $1/32$ .

1. Calculate the entropy of this source,
2. Calculate the entropy of a source that produces eight equally probable symbols, and
3. Compare the two entropies.

# Entropy Theory - Reminder

- Quantifies the average amount of information of a source,
- Average number of bits of information required to represent the symbols the source produces

$$H(X) = \sum_{i=1}^n P(x_i) I(x_i) = - \sum_{i=1}^n P(x_i) \log P(x_i)$$

- X represents the symbols of an alphabet.
- "A mathematical theory of communication,"
  - C. E. Shannon, *Bell Systems Technical Journal*, 1948
- Refer to the lecture slides for more information. (04 - Information Theory)

# Entropy Calculation Exercise – Question 1

Symbol	$x_0$	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$
Probability	1/4	1/4	1/8	1/8	1/8	1/16	1/32	1/32

For  $x_0$  and  $x_1$ , we have:  $-1/4 \log_2 1/4 = -1/4 \log_2 4^{-1} = 1/4 \log_2 2^2 = 2/4$ , since  $\log_2 2 = 1$

For  $x_2$ ,  $x_3$  and  $x_4$ , we have:  $-1/8 \log_2 1/8 = -1/8 \log_2 8^{-1} = 1/8 \log_2 2^3 = 3/8$ , since  $\log_2 2 = 1$

For  $x_5$  we have:  $-1/16 \log_2 1/16 = -1/16 \log_2 16^{-1} = 1/16 \log_2 2^4 = 4/16$ , since  $\log_2 2 = 1$

For  $x_6$  and  $x_7$ , we have:  $-1/32 \log_2 1/32 = -1/32 \log_2 32^{-1} = 1/32 \log_2 2^5 = 5/32$ , since  $\log_2 2 = 1$

$$H(X) = \sum_{i=1}^n P(x_i) I(x_i) = - \sum_{i=1}^n P(x_i) \log P(x_i)$$

Thus, we have:

$$\begin{aligned} H(X) &= 2/4 + 2/4 + 3/8 + 3/8 + 3/8 + 4/16 + 5/32 + 5/32 \\ &= (16 + 16 + 12 + 12 + 12 + 8 + 5 + 5) / 32 = \\ &= 86/32 = 2.6875 \end{aligned}$$

# Entropy Calculation Exercise – Question 2

We now have 8 symbols with equal probability, thus the probability for each one of them is 1/8.

Symbol	$x_0$	$x_1$	$x_2$	$x_3$	$x_4$	$x_5$	$x_6$	$x_7$
Probability	1/8	1/8	1/8	1/8	1/8	1/8	1/8	1/8

$$H(X) = \sum_{i=1}^n P(x_i) I(x_i) = - \sum_{i=1}^n P(x_i) \log P(x_i) \longrightarrow H(X) = - \sum_{i=1}^n \frac{1}{n} \log \frac{1}{n} = \log n$$

In this question we have a source with random behavior, thus we can conclude that it requires  $\log_2 8 = \log_2 2^3 = 3$ , since  $\log_2 2 = 1$

# Entropy Calculation Exercise – Question 3

- Obviously,  $3 > 2.6875$ , but *what conclusion can we reach about random sources?*
- Q1: some symbols have higher probabilities, meaning they contribute less to the overall uncertainty of the source.
  - The entropy is lower, indicating less unpredictability in the information produced by this source.
- Q2: there is maximum uncertainty about which symbol will be produced, as no symbol is favored over another.

Source	Q1	Q2
H(X)	2.6875	3

- The more evenly distributed the probabilities across possible outcomes, the higher the entropy.
- Random sources maximize entropy.
- Higher entropy values indicate greater uncertainty or unpredictability in the source's output.
- Lower entropy suggests a more predictable source.